

(Im)Possibilities of Predicate Detection in Crash-Affected Systems using Interrupt-Style Failure Detectors

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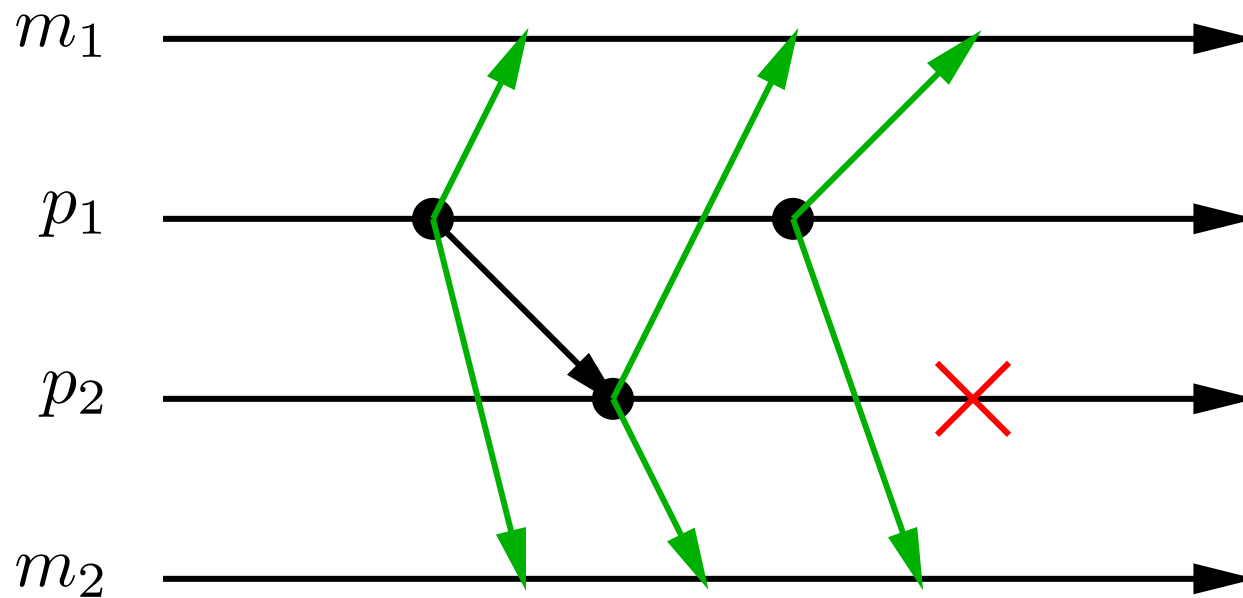
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Predicate Detection



- Does a global predicate φ hold throughout the computation?

Predicate Detection Semantics

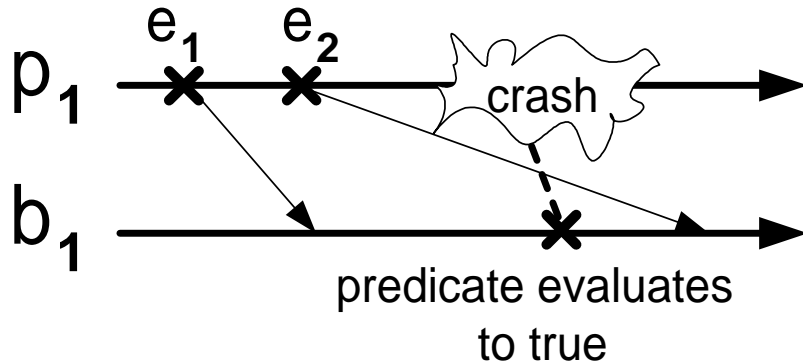
- **Perfect predicate detection** Sem_1 :
 - (S) If the algorithm triggers a detection, then φ has held in the computation.
 - (L) If φ holds, then the algorithm will eventually trigger a detection.
- **Stabilizing predicate detection** Sem_2 :
 - L and $\diamond S$.
- **Infinitely often accurate predicate detection** Sem_3 :
 - L and $\square \diamond S$.

Focus

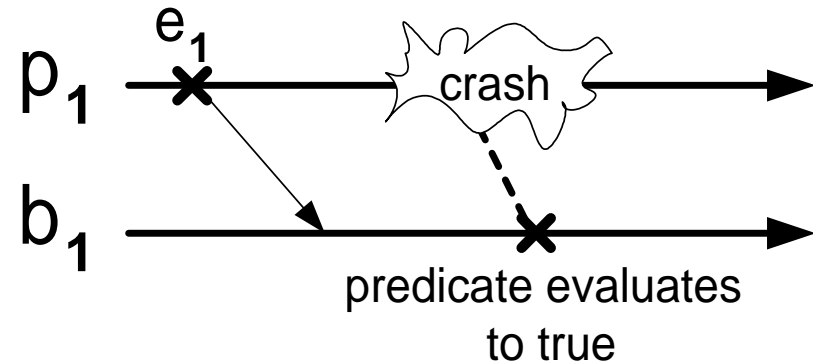
- Which **predicate detection semantics** are achievable in asynchronous systems where **crash faults** can happen?
- Use asynchronous systems augmented with **(unreliable) failure detectors**.
- Relevant failure detector classes:
 - **Perfect** \mathcal{P} and **eventually perfect** $\diamond\mathcal{P}$ [CT96].
 - **Infinitely often accurate** $\square\diamond\mathcal{P}$ [GM98].

Most interesting Result

- \mathcal{P} **not** sufficient for perfect predicate detection.



(a)



(b)

Types of Failure Detectors

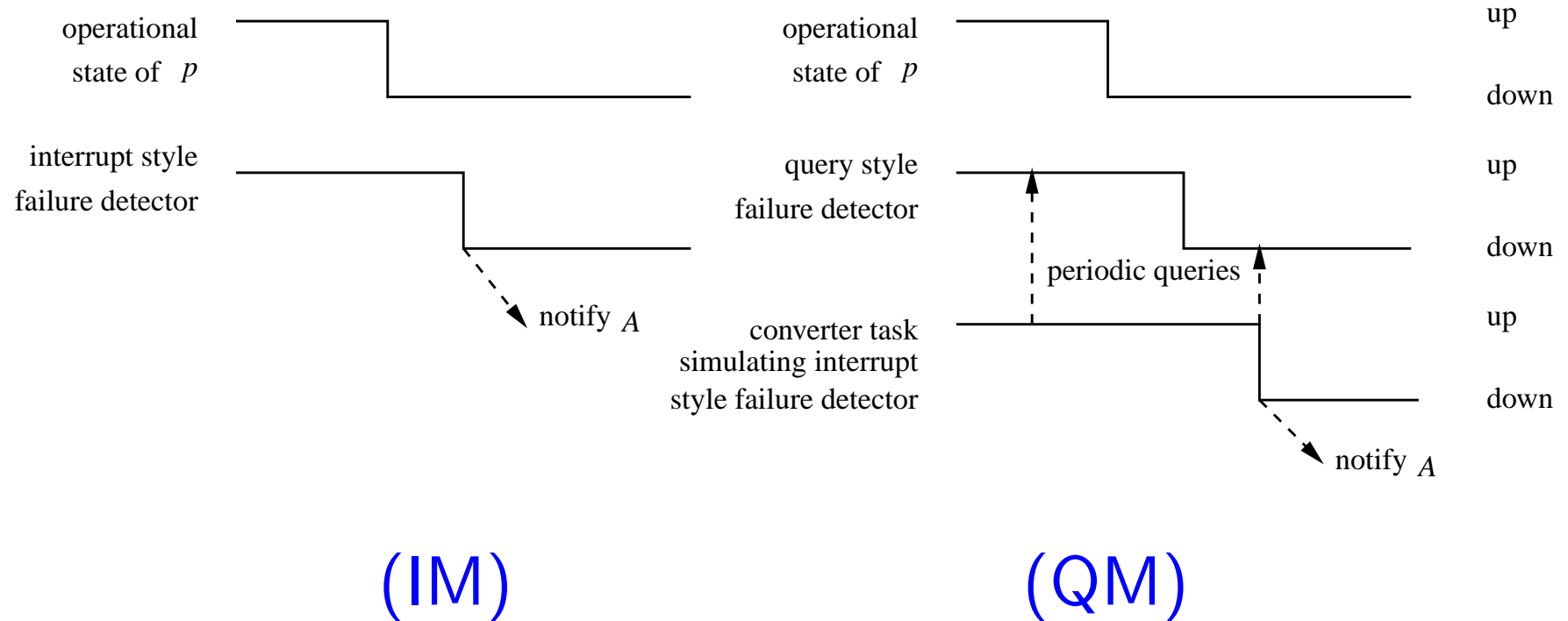
- **QM – Query model** (Chandra and Toueg [CT96]):
Query-style failure detectors.
 - Spurious detections can go unnoticed.
- **IM – Interrupt model** (Garg and Mitchell [GM98]):
Interrupt-style failure detectors.
 - Every detection reaches application.
- We use interrupt-style ones.

Solvability of Problems

- If a problem P is solvable using \mathcal{D} in QM, then P is solvable in IM.
 - Proof idea: IM is more restrictive.
- **Vice versa?** (P solvable in IM \Rightarrow P solvable in QM.) ■
- Only for $\mathcal{D} \in \mathcal{P}$ or $\diamond\mathcal{P}$, not for $\mathcal{D} \in \square\diamond\mathcal{P}$.

Proof Idea

- Use **converter task** which regularly queries failure detector.



Summary

- Predicate detection in crash-affected systems.
- Which predicate detection semantics are achievable using which types of failure detectors?
- Must go for stabilizing predicate detection semantics in many practical settings.
- Interesting aspect: IM vs. QM.
- For more details see WSS paper [GP01a] and IBM Research Report [GP01b] for full proofs.

Acknowledgements

- Slides produced using pdfL^AT_EX and Klaus Guntermann's PPower4.

References

- [CT96] Tushar Deepak Chandra and Sam Toueg. Unreliable failure detectors for reliable distributed systems. *Journal of the ACM*, 43(2):225–267, March 1996.
- [GM98] Vijay K. Garg and J. Roger Mitchell. Distributed predicate detection in a faulty environment. In *Proceedings of the 18th IEEE International Conference on Distributed Computing Systems (ICDCS98)*, 1998.
- [GP01a] Felix C. Gärtner and Stefan Pleisch. (Im)Possibilities of predicate detection in crash-affected systems. In *Proceedings of the 5th Workshop*

on Self-Stabilizing Systems (WSS 2001), Lecture Notes in Computer Science, Lisbon, Portugal, October 2001. Springer-Verlag. to appear.

[GP01b] Felix C. Gärtner and Stefan Pleisch. (Im)Possibilities of predicate detection in crash-affected systems. Research Report RZ 3361 (#93407), IBM Research Laboratory, Zurich, August 2001.